The Noise Collector for sparse recovery in high dimensions

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Inverse problems in wave propagation



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Imaging setup



An signal is emitted from x_s at the array of N transducers, it illuminates the scatterers in the image window (IW). The point scatterers are at y_j , and now they can be viewed as the sources of the signal. They send the scattered signal back to the array. There are K pixels in the IW. The number of scatters is M < N, and N < K, typically. The map $A\rho \Rightarrow b$ in the *paraxial approximation* is (up to a constant) the (partial) Fourier transform.

Imaging of sparse scenes



Left: the true image. I show 2-dimensional images for simplicity. Right: the recovered solution vector is plotted with red stars and the true solution vector of $A\rho = b$ with green circles.

Noise, l_1 versus l_2 regularization



 l_1 -methods are unstable to noise., l_2 -methods loose resolution.

l_1 -regularization and Lasso

Noiseless. Want to solve a sparsity promoting optimization

$$\rho = \arg \min \|\tilde{\rho}\|_0$$
, subject to $A\tilde{\rho} = b$

where $\|\rho\|_0 = \{\#\rho_i \neq 0\}$. It is expensive, so we solve

$$\rho = \arg \min \|\tilde{\rho}\|_1$$
, subject to $A\tilde{\rho} = b$

where $\|\rho\|_1 = \sum_i |\rho_i|$. Noisy case, Lasso. R. Tibshirani '96, Chen & D.Donoho '94, F.Santosa & W.Symes '86

$$\rho = \arg\min\left(\lambda \|\tilde{\rho}\|_1 + \frac{\|A\tilde{\rho} - b\|_2^2}{2}\right)$$

where $\|\rho\|_2 = \sqrt{\sum_i |\rho_i|^2}$ and λ is a tuning parameter.

Tuning λ in Lasso



LASSO results with $\lambda = 1$, $\lambda = 0.5$ (optimal) and $\lambda = 0.1$.

l_1 & LASSO

LASSO: finds a *sparse* approximate solution $A\rho \approx b$ if the tuning parameter λ is chosen correctly.

 l_1 : finds a *sparse* solution only if noise e = 0. No tuning parameters. If e is small, the sparse approximate solution can be found by thresholding. Thresholding has to be tuned.

We propose to solve

 $(\rho_{\tau},\eta) = \arg\min\left(\tau \|\rho\|_{1} + \|\eta\|_{1}\right)$, subject to $A\rho + C\eta = b$

where C is the noise collector matrix and τ is the weight of the noise collector, and $b = b_0 + e$. We can prove $\rho \approx \rho_{\tau}$ and $C\eta \approx e$ if τ is chosen correctly. We can choose $\tau = O(\sqrt{\ln N})$ for any level of noise, before de-noising.



The top row - the images, the bottom row - the solution vector with red stars and the true solution vector with green circles.



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Design of the Noise Collector

- (i) Columns of NC should be sufficiently orthogonal to the columns of A, so it does not absorb signals with meaningful information.
- (ii) Columns of NC should be uniformly distributed on the unit sphere \mathbb{S}^{N-1} so that we could approximate well a typical noise vector.
- (iii) The number of columns of NC should grow slower than exponential with N, otherwise the method is impractical.
- (iv) Deterministic approach. If we fill up C imposing

$$|\vec{a}_i \cdot \vec{c}_j| < \frac{\alpha}{\sqrt{N}} \,\forall i, j \,, \text{ and } |\vec{c}_i \cdot \vec{c}_j| < \frac{\alpha}{\sqrt{N}} \,\forall i \neq j,$$
 (1)

then the Kabatjanskii-Levenstein inequality implies that the number Σ of columns in C grows at most polynomially: $N^{\alpha} \leq \Sigma \leq N^{\alpha^2}$.

Probabilistic approach to design of NC

If the columns of C are drawn at random independently. then the dot product of any two random unit vectors is still typically of order $1/\sqrt{N}$. We have an asymptotically negligible event that our noise collector is bad The decoherence constraint is weakened by a logarithmic factor.

Lemma: Choose $\beta > 1$, and pick $\Sigma = N^{\beta}$ vectors $\vec{c_i}$ at random and independently on \mathbb{S}^{N-1} . Then, for any $\kappa > 0$ there are constants $c_0(\kappa, \beta)$ and $\alpha > 1/2$, such that (i)

$$|\vec{a}_i \cdot \vec{c}_j| < c_0 \sqrt{\ln N} / \sqrt{N} \quad \text{for all } i, j, \tag{2}$$

and (ii) for any $\vec{e} \in \mathbb{S}^{N-1}$ there exists at least one \vec{c}_j , so

$$|\vec{e} \cdot \vec{c}_j| > \alpha/\sqrt{N},\tag{3}$$

with large probability $1 - 1/N^{\kappa}$. In addition the condition number of [A|C] is O(1).

False Discovery Rate is zero

Theorem 1: (No phantom signal) Suppose there is no signal: $\rho = 0$ and $e/||e||_{l_2}$ is uniformly distributed on the unit sphere. For any $\kappa > 0$ we can construct the noise collector and choose weight τ so that $\rho_{\tau} = 0$ with probability $1 - 1/N^{\kappa}$.

Theorem 2: Let ρ be an *M*-sparse solution of $A\rho = b_0$. If the columns of A are decoherent: $|a_i \cdot a_j| \leq \frac{1}{3M}$, then $\operatorname{supp}(\rho_{\tau}) \subseteq \operatorname{supp}(\rho)$ with probability $1 - 1/N^{\kappa}$.

Supports of ρ and ρ_{τ} agree

Theorem 3: Suppose r is the magnitude of smallest non-zero entry of ρ . If $||e||_{l_2}/||b_0||_{l_2} \leq c_2 \sqrt{\ln N}$, $c_2 = c_2(\kappa, \beta, r, M)$, then $\operatorname{supp}(\rho_{\tau}) = \operatorname{supp}(\rho)$, with probability $1 - 1/N^{\kappa}$.

Theorem 4: (Exact Recovery): If there is no noise e = 0. Then $\rho_{\tau} = \rho$ with probability $1 - 1/N^{\kappa}$.

Failure to recover



Image Window and Noise Collector



GeLMA with Fast Noise Collector

Require: Set
$$\rho = 0$$
, $z = 0$. $\eta = 0$.
Pick $\beta = \beta(A, C)$, and $\tau = .8\sqrt{\ln N}$.
repeat
 $r = b - A\rho - C\eta$

$$\rho \Leftarrow \mathcal{S}_{\tau\beta}(\rho + \beta A^*(z+r))$$

$$\eta \Leftarrow \mathcal{S}_{\beta}(\eta + \beta C^*(z+r))$$

until Convergence

Calibrate τ so that FDR is zero when b = e (only noise). "No phantom signal" criterion.

Fast Noise Collector C is several random circular matrices. Then it is cheap to store C, and we can use FFT for matrix-vector multiplication.

The soft shrinkage-thresholding operator

$$\mathcal{S}_{\tau}(y_i) = \operatorname{sign}(y_i) \max\{0, |y_i| - \tau\}.$$

Geometric interpretation of \vec{z} $\vec{a}_i \cdot \vec{z} = \tau \operatorname{sign}(\rho_i)$, if $\rho_i \neq 0$, and $|\vec{a}_i \cdot \vec{z}| \leq \tau$ if $\rho_i = 0$. Assume both \vec{a}_i and $-\vec{a}_i$ are columns of A, and all $\rho_i \ge 0$ then

$$H_A = \{ x \in \mathbb{R}^N : x = \sum_i \alpha_i \vec{a}_i, \sum_i \alpha_i \leqslant 1, \alpha_i \geqslant 0 \}$$

Suppose Λ is the support of ρ , typically (for non-sparse ρ) $|\Lambda| = N$. Then, the simplex

$$\left\{ \vec{x} \in \mathbb{R}^N \left| \vec{x} = \sum_{i \in \Lambda} \alpha_i \vec{a}_i, \sum_{i \in \Lambda} \alpha_i = 1, \alpha_i \ge 0 \right. \right\}$$

has the unique normal vector \vec{n} , which is collinear to \vec{z} because

$$1 = \vec{z} \cdot \vec{a}_i = \frac{\vec{z} \cdot \vec{b}}{\|\vec{b}\|_A}, \ \forall i \in \Lambda, \text{ and } \vec{z} \cdot \vec{a}_j < 1, \forall j \notin \Lambda.$$
(4)

Geometric interpretation of z



If A is the identity matrix then $\Phi(e) = (\operatorname{sign}(e_1), \ldots, \operatorname{sign}(e_N)).$

Duality

Given A we have H_A define $\vec{z} = \Phi_A(\vec{b})$. Let $Z = \{ \text{ all } \vec{z} = \Phi_A(\vec{e}) \text{ for some } \vec{e} \in \mathbb{S}^{N-1} \}$

Then

$$\max_{z \in Z} z \cdot b = \min(\|\rho\|_1, \text{ subject to } A\rho = b),$$

and

$$\Phi_A(\vec{b}) = \arg \max_{z \in Z} z \cdot b.$$

Proof of Theorem 1

Solve

$$(\rho_{\tau},\eta) = \arg\min\left(\tau\|\rho\|_1 + \|\eta\|_1\right)$$
, subject to $A\rho + C\eta = b$

Theorem 1:(No phantom signal) Suppose there is no signal: $\rho = 0$ and $e/||e||_{l_2}$ is uniformly distributed on the unit sphere. For any $\kappa > 0$ we can construct the noise collector and choose weight τ so that $\rho_{\tau} = 0$ with probability $1 - 1/N^{\kappa}$.

Proof: Instead of Φ_A , consider $\Phi_C : e \to z$, where z is dual certificate of optimality of η . We want to show $\Phi_{[\tau A|C]} : e \to z$. It means that z is the also dual certificate of optimality of $(0, \eta)$, i.e.

$$|\vec{z} \cdot \vec{a}_j| < \tau, \forall j$$

 $\operatorname{Map}\, \Phi_C: e \to z$



Random vectors on \mathbb{S}^{N-1}

Everything is rotation invariant in $\Phi_C : e \to z$. Thus $n = z/||z||_2$ is uniformly distributed on \mathbb{S}^{N-1} . $||z|| = O(\sqrt{N})$, because l_1 -balls are probabilistically l_2 -small. Coordinates of s uniformly distributed vector on \mathbb{S}^{N-1} have i.i.d. Gaussian distribution $N(0, 1/\sqrt{N})$. The event $|\vec{z} \cdot \vec{a}_j| < \tau, \forall j$ does not happen with probability

$$\mathbb{P}\left(\max_{j} |\vec{z} \cdot \vec{a}_{j}| \ge \tau\right) \leqslant N^{\beta} \mathbb{P}\left(|\vec{n} \cdot \vec{a}_{1}| \ge \tau/\sqrt{N}\right)$$
$$\leqslant 2N^{\beta} \mathrm{e}^{-c\tau^{2}} \leqslant 1/N^{\kappa}, \text{ if } \tau = O(\sqrt{\ln N}).$$